

# Methods of Proof for Boolean Logic

## Proof by Contradiction

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## Outline

- 1 Proof by Contradiction
- 2 Arguments With Inconsistent Premises

## Announcements

For 02.17

- 1 HW4 is due now
- 2 HW1-3 will be returned on Thursday

## Proof by Contradiction

### Proving a Negated Claim

- Okay, we've done proofs for conjunction and disjunction, but what about negation?
- How would you go about proving a negated claim like  $\neg \text{Cube}(a)$ ?
- Well,  $\neg \text{Cube}(a)$  is true if and only if  $\text{Cube}(a)$  is false
- So, we can prove  $\neg \text{Cube}(a)$  by showing that  $\text{Cube}(a)$  is false
- But how do we do that?

# Proof by Contradiction

Proving Something False

- There is a very important method for proving something false
  - **Proof by Contradiction**
  - A.k.a **Indirect Proof, Reductio ad Absurdum**

## The Basic Idea of Proof by Contradiction

To show that P is false, it suffices to show that something which cannot possibly be true, i.e. a **contradiction**, follows from P

# Proof by Contradiction

What is a *Contradiction*?

## Contradiction

- Intuitively, a **contradiction** is any sentence that cannot possibly be true, or any group of sentences that cannot all be true **simultaneously**
- The symbol  $\perp$  is often used as a short-hand way of saying that a contradiction has been obtained

## Examples:

- 1  $\neg\text{Cube}(a) \wedge \text{Cube}(a)$
- 2  $a = b, b = c, a \neq c$

# Proof by Contradiction

What is It?

## Proof by Contradiction (Preliminary Version)

To prove that P is false, show that a contradiction  $\perp$  follows from P

## Proving a Negated Claim

To prove  $\neg P$ , assume P and prove a contradiction  $\perp$

- Contradictions are **impossible**, so **false**
- If you can show that P leads to a contradiction, then P must be false
- But if P is false, then  $\neg P$  must be true

# Proof by Contradiction

A Simple Example

**Claim:** This argument is valid

$$\begin{array}{l} \neg\text{SameShape}(a, b) \\ b = c \\ \hline \neg a = c \end{array}$$

**Proof:** We want to show  $\neg a = c$  from the premises, so we will use a **proof by contradiction**

- 1 Suppose  $a = c$
- 2 Then, from **premise one**  $\neg\text{SameShape}(c, b)$  follows by Indiscernibility of Identicals
- 3 But by **premise two**, we know  $\text{SameShape}(c, b)$ . **This** is a contradiction,  $\perp$ !
- 4 So our **supposition** must have been **false**; that is,  $\neg a = c$  must be **true** given the premises

# Proof by Contradiction

In Class Exercise

Write an informal proof of this argument. Do a proof by contradiction.

$$\text{Cube}(b) \vee \text{LeftOf}(b, c)$$

$$\text{SameCol}(b, c)$$

$$\neg \text{Tet}(b)$$

# Proof by Contradiction

Official Version

## Proof by Contradiction (Official Version)

- 1 To prove that P is false, show that a contradiction  $\perp$  follows from P
  - 2 To prove that P is true, show that a contradiction  $\perp$  follows from  $\neg P$
- Proof by contradiction can also be used for proving **un-negated claims**

# Proof by Contradiction

Used to Prove an Un-negated Claim

$$b = c \wedge \text{SameShape}(c, a)$$

$$\text{Cube}(b)$$

$$\text{Cube}(a)$$

**Proof:** We will use a proof by contradiction. First, suppose the conclusion is false, i.e.  $\neg \text{Cube}(a)$ . By premise 1 we know that  $\text{SameShape}(c, a)$ , so it must be that  $\neg \text{Cube}(c)$ . But premise 1 tells us that  $b = c$  and together with premise 2 this entails that  $\text{Cube}(c)$ , by the Indiscernibility of Identicals. This contradicts our early conclusion, so our supposition must have been false. Thus, the conclusion must be true when the premises are.

# Proof by Contradiction

In Class Exercise

Write an informal proof of this argument, phrased in complete, well-formed English sentences. **Hint:** try a proof by contradiction. I encourage you to work in groups.

5.15

$$\text{Tet}(b)$$

$$\text{Cube}(c)$$

$$\text{Larger}(c, b) \vee c = b$$

$$\text{Smaller}(b, c)$$

# Proof by Contradiction

## A More Advanced Example

### Proof

We will do a proof by cases based on **premise three**:

- 1 Suppose  $\neg \text{Small}(d)$ . We want to show  $d = c \vee d = b$ , let's do this by indirect proof. This requires deriving  $\perp$  from the additional supposition that  $\neg(d = c \vee d = b)$ . First, note that by DeMorgan's Laws this implies  $d \neq c$  and  $d \neq b$ . From this and **our first supposition**, **premise two** clearly requires that  $\neg \text{Dodec}(d)$ . But, from **premise four** and **our original supposition**, **premise one** clearly requires that  $\text{Dodec}(d)$ . **These requirements** are contradictory,  $\perp$ .

- 2 Now suppose  $d = b$ . Then  $d = c \vee d = b$  follows immediately by disjunction intro

In either case  $d = c \vee d = b$  follows, so the argument is valid

### Argument 4

$\text{Dodec}(d) \vee \text{Tet}(d) \vee \text{Small}(d)$   
 $\neg \text{Dodec}(d) \vee (d = c) \vee \text{Small}(d)$   
 $\neg \text{Small}(d) \vee d = b$   
 $\neg \text{Tet}(d)$   
 $d = c \vee d = b$

- We had a proof by contradiction **inside** a proof by cases!
- This is like exercise 5.17!

# Valid Arguments

## What If The Premises are Inconsistent?

### Logical Consequence, Validity

$C$  is a logical consequence of  $P_1, \dots, P_n$  if and only if it is **impossible** for  $P_1, \dots, P_n$  to be true while  $C$  false

- Now think about an argument with inconsistent premises
  - Is it valid?
- Yes!
  - Why: it's impossible for the premises to be true
- So, it is impossible for the premises to be true while the conclusion is false!
- But, crucially, the argument is **not sound**

# Summary

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- We learned a powerful proof method:
  - Proof by contradiction
- We wrapped our head around the fact that any argument with contradictory premises is valid
  - Importantly, though, **such an argument is never sound**