

# Expressing *May* and *Must*

## Dynamic Semantics at Work

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# The Role of Information

## In Inquiry and Conversation

- Informational contents (*propositions*) are sets of possible worlds
  - These sets distinguish ways world might be (worlds in the set) from ways it isn't (worlds excluded from set)
- One informational content is particularly useful for understanding how linguistic interactions unfold:

### Contextual Possibilities (*c*)

As communication and inquiry unfold, a body of information accumulates. Think of this information as what the agents are mutually taking for granted in some way. I call the set of worlds embodying this information *c*, short for *contextual possibilities*. (Stalnaker 1978; Lewis 1979)

# Outline

- 1 Free Choice and Two Views of Semantics
- 2 Dynamic Preference Semantics

# Gaining Information

## And Eliminating Possibilities

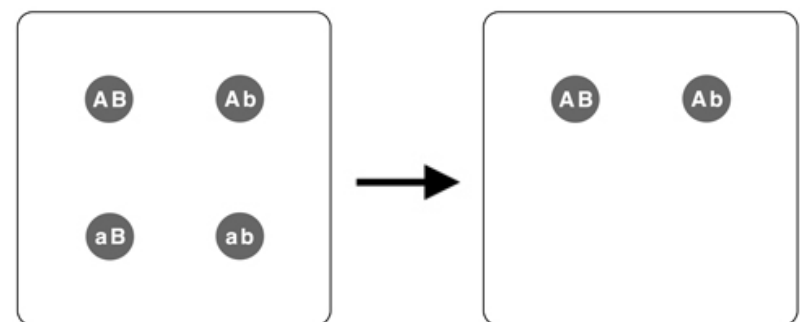


Figure: Accepting the information that A

- Inquiry progresses by using information to eliminate uncertainty, i.e. the elimination of worlds.
- $\{w_{AB}, w_{Ab}, w_{aB}, w_{ab}\} \Rightarrow \{w_{AB}, w_{Ab}\}$

## Question

How are Semantics and Information Change Related?

### Classical Picture

- 1 **Semantics:** pair sentences w/propositions
  - $\llbracket \phi \rrbracket$  is a set of worlds
- 2 **Pragmatics:** rules for rational agents
  - When presented with information, rational agents use it to eliminate possibilities (update using intersection)

The Point **Semantics** specifies informational content of a sentence, but nothing in particular about how sentence changes contextual information ( $c$ )

- Instead, **pragmatics** says how the sentence's content impacts  $c$ :  $c \cap \llbracket \phi \rrbracket$

## The Classical Picture

In More Detail

### Classical Possible Worlds Semantics

- 1  $\llbracket A \rrbracket = \{w \mid w(A) = 1\}$
- 2  $\llbracket \neg \phi \rrbracket = W - \llbracket \phi \rrbracket$
- 3  $\llbracket \phi \wedge \psi \rrbracket = \llbracket \phi \rrbracket \cap \llbracket \psi \rrbracket$
- 4  $\llbracket \phi \vee \psi \rrbracket = \llbracket \phi \rrbracket \cup \llbracket \psi \rrbracket$
- 5  $\llbracket \Diamond \phi \rrbracket = \{w \mid \exists w' : w' \in Acc(w) \ \& \ w' \in \llbracket \phi \rrbracket\}$ 
  - $Acc(w)$  is the set of worlds accessible from  $w$

### Classical Truth and Consequence

**Truth**  $w \models \phi \iff w \in \llbracket \phi \rrbracket$

**Consequence**  $\phi \models \psi \iff \llbracket \phi \rrbracket \subseteq \llbracket \psi \rrbracket$

## Two Important Consequences

Of the Classical Approach

### Fact 1: $\Diamond A \models \Diamond A \vee \Diamond B$

- Consequence is content inclusion:
  - $\llbracket \Diamond A \rrbracket \subseteq \llbracket \Diamond A \vee \Diamond B \rrbracket$
- Disjunction is union:
  - $\llbracket \Diamond A \rrbracket \subseteq \llbracket \Diamond A \rrbracket \cup \llbracket \Diamond B \rrbracket$

### Fact 2: $\Diamond A \vee \Diamond B \not\models \Diamond A$

- This would require:
  - $\llbracket \Diamond A \rrbracket \cup \llbracket \Diamond B \rrbracket \subseteq \llbracket \Diamond A \rrbracket$
  - But this only holds when  $\llbracket \Diamond B \rrbracket = \emptyset$

## The First Problem of Free Choice

Kamp (1973) and Ross (1941)

(1) Billy **may** go to the beach

- Best classical approximation of *May*: some species of  $\Diamond$
- $May B \stackrel{?}{=} \Diamond B$

**Problem:**  $May B$  does not entail  $May B \vee May C$

X: Billy may go to the beach

Y: Ah, so Billy may go to the beach, or Billy may go the cinema  $\times \times$

**Why is this a problem?**

- Fact 1:  $\Diamond B \models \Diamond B \vee \Diamond C$

## The Second Problem of Free Choice

Kamp (1973)

(2) Billy may go to the beach or he may go to the cinema

Problem:  $\text{May } B \vee \text{May } C$  **does** entail  $\text{May } B$

X: Billy may go to the beach or he may go to the cinema

Y: Ah, so Billy may go to the beach ✓

- Indeed:  $\text{May } B \vee \text{May } C$  entails  $\text{May } B \wedge \text{May } C$

Why is this a problem?

- Fact 2:  $\diamond A \vee \diamond B \neq \diamond A$
- So there are at least two problems for classical approach to *may*

## Tentative Diagnosis

The Performative Function of *May*

- A pragmatic approach...
  - Competence with a **word** is involved, so a semantic approach is preferable
  - Also, classical pragmatics assumes (declarative) sentences denote propositions
    - Proposition gets intersected with  $c$
  - $\text{May } A$  shouldn't change  $c$ , but rather  $Acc$
  - This operation can't be intersection (sphere of accessibility should **grow** w/permission)
  - This would also require  $\text{May } A$  to denote an accessibility relation
    - So many problems w/this
- Alternative: approach to semantics that can directly encode (features of) communicative function of a word

## Question

How are Semantics and Information Change Related?

Classical Picture

- 1 **Semantics:** pair sentences w/propositions
  - $[[\phi]]$  is a set of worlds
- 2 **Pragmatics:** rules for rational agents
  - When presented with information, rational agents use it to eliminate possibilities (update using intersection)

Dynamic Picture (Veltman, Heim)

- 1 **Semantics:** pair sentences with ways of changing contexts
  - $\phi$  eliminates worlds from  $c$
- 2 **Pragmatics:** mechanisms for additional changes
  - General rationality, theory of mind, etc.

## The Dynamic Picture

In More Detail

The Basic Idea

Assign each  $\phi$  a function  $[\phi]$  encoding how it changes  $c$ :  
 $c[\phi] = c'$  (Better notation:  $[\phi](c) = c'$ )

Dynamic Informational Semantics

- 1  $c[A] = \{w \in c \mid w(A) = 1\}$
- 2  $c[\neg\phi] = c - c[\phi]$
- 3  $c[\phi \wedge \psi] = (c[\phi])[ \psi ]$
- 4  $c[\phi \vee \psi] = c[\phi] \cup c[\psi]$
- 5  $c[\diamond\phi] = \{w \mid \exists w' : w' \in Acc(w) \ \& \ w' \in c[\phi]\}$ 
  - Veltman (1996): for epistemic *might*,  $Acc(w)$  is  $c$

## Dynamic Informational Semantics

Semantic Concepts and a New Perspective on Classical Ones

### Support, Consequence (Veltman)

- $c \models \phi \iff c[\phi] = c$
- $\phi_1, \dots, \phi_n \models \psi \iff \forall c : c[\phi_1] \dots [\phi_n] \models \psi$

### Truth, Propositions (Starr)

$$w \models \phi \iff \{w\}[\phi] = \{w\} \quad \llbracket \phi \rrbracket = \{w \mid w \models \phi\}$$

### Classical Consequence (Starr)

$$\phi_1, \dots, \phi_n \models \psi \iff \forall w : \{w\}[\phi_1] \dots [\phi_n] \models \psi$$

- Classical logic is the logic of perfect information
- Equivalent to standard  $\models$  for modal free fragment

## Dynamic Informational Semantics

Understanding Classical vs. Dynamic Semantics

### Corollary (Boolean Equivalence)

- 1  $\llbracket A \rrbracket = \{w \mid w(A) = 1\}$
- 2  $\llbracket \neg \phi \rrbracket = W - \llbracket \phi \rrbracket$
- 3  $\llbracket \phi \wedge \psi \rrbracket = \llbracket \phi \rrbracket \cap \llbracket \psi \rrbracket$
- 4  $\llbracket \phi \vee \psi \rrbracket = \llbracket \phi \rrbracket \cup \llbracket \psi \rrbracket$

- Booleans behave classically
- Boolean fragment:  $\models$  and  $\models$  are identical
  - It's only with modals that  $\models$  and  $\models$  come apart
- **Surprise:**  $w \notin \llbracket \diamond \phi \rrbracket$  unless  $Acc(w) = \{w\}$  and  $w \in \llbracket \phi \rrbracket$
- Then  $\llbracket \diamond \phi \rrbracket \subseteq \llbracket \phi \rrbracket$  and so  $\diamond \phi \models \phi$ 
  - Yet  $\diamond \phi \not\models \phi$

## Illustrating Our Surprise

$\diamond A$  is false unless  $A$  is true and  $Acc(w) = \{w\}$

- Recall  $w \in \llbracket \diamond A \rrbracket \iff \{w\}[\diamond A] = \{w\}$
- Consider whether  $w_A \in \llbracket \diamond A \rrbracket$ :

$$\begin{aligned} \{w_A\}[\diamond A] &= \{w \mid \exists w' : w' \in Acc(w) \ \& \ w' \in \{w_A\}[A]\} \\ &= \{w \mid w_A \in Acc(w)\} \end{aligned}$$

- Unless  $Acc(w_A) = \{w_A\}$ ,  $\{w_A\}[\diamond A] \neq \{w_A\}$
- Consider whether  $w_a \in \llbracket \diamond A \rrbracket$

$$\begin{aligned} \{w_a\}[\diamond A] &= \{w \mid \exists w' : w' \in Acc(w) \ \& \ w' \in \{w_a\}[A]\} \\ &= \{w \mid \exists w' : w' \in Acc(w) \ \& \ w' \in \emptyset\} \\ &= \emptyset \end{aligned}$$

- Nope! So  $\llbracket \diamond A \rrbracket \subseteq \llbracket A \rrbracket$

## Classical Logic as Logic of Omniscience

A Logic of Omniscience is Not Suited for Modality

### d'Alembert (1751) on Truth

“The universe... would only be one fact and one great truth for whoever knew how to embrace it from a single point of view.” (d'Alembert 1995:29)

- A sentence is true iff it corresponds w/‘the great truth’
  - Correspond: cohere w/the great truth,  $\{w\}[\phi] = \{w\}$
- But modal sentences are indexical to **uncertainty** or **utilities** of agent(s) evaluating them
- Whether they cohere with the great truth is irrelevant
- We want an evaluative concept that targets how they cohere with **partial information**:  $c[\phi] = c$  (support)

# Dynamic Informational Semantics

Back to Free Choice

- This system is quite useful for **epistemic** modals
  - $A \wedge \Diamond \neg A$  is inconsistent (cf. Veltman and Yalcin)
  - $\Diamond \phi$  can add possibilities to  $c$ 
    - Not possible to get with  $c \cap \llbracket \Diamond \phi \rrbracket$
- But it does **not** solve the problem of free choice:
  - $\Diamond \phi \models \Diamond \phi \vee \Diamond \psi$
  - $\Diamond \phi \vee \Diamond \psi \not\models \Diamond \phi$
- However, it is the backbone of a system that does:
  - *may*'s meaning resides in how it changes context
  - Use of  $\models$  rather than  $\vDash$

## Basic Idea

Capture the non-informational impact of *may* by sophisticating the contexts sentences operate on

# Preference, Rationality & Context

Information and the Process of Inquiry

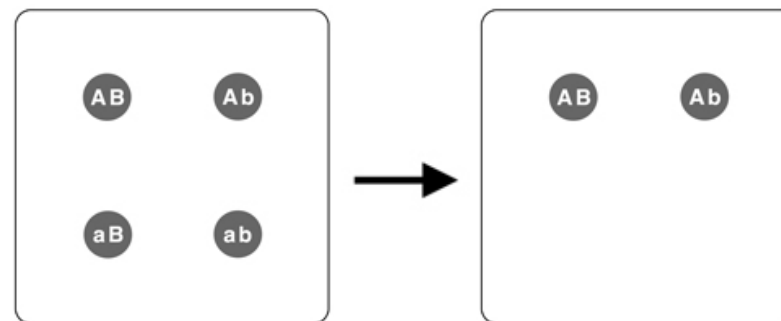


Figure: Accepting the information that A

- Inquiry progresses by gaining information, i.e. the elimination of worlds.
- $\{w_{AB}, w_{Ab}, w_{aB}, w_{ab}\} \Rightarrow \{w_{AB}, w_{Ab}\}$

# Preference, Rationality & Context

Issues

- It's not just information that accumulates in communication and inquiry (Bromberger 1966)
- There are issues (e.g. Hamblin 1958; Roberts 1996).
- They can be thought of as ways of grouping worlds in  $c$  into competing alternative propositions.

## Alternatives ( $C$ ) (e.g. Groenendijk 1999)

Alternatives represent open, competing propositions the agents are concerned with deciding between; their **issues**. Formally, this grouping of  $c$  may be identified with a set of sets of worlds; call it  $C$ . There is no need to also keep track of  $c$ : it is just the union of all the alternatives in  $C$ .

# Preference, Rationality & Context

Issues and Inquiry

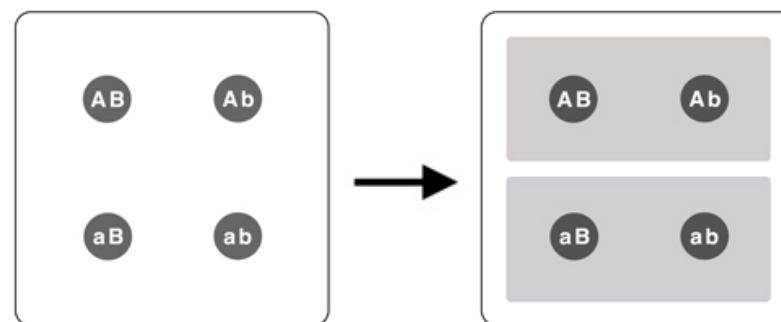


Figure: Recognizing the issue whether A

- Inquiry also progresses by recognizing issues, i.e. introducing alternatives
- $\{\{w_{AB}, w_{Ab}, w_{aB}, w_{ab}\}\} \Rightarrow \{\{w_{AB}, w_{Ab}\}, \{w_{aB}, w_{ab}\}\}$

# Preference, Rationality & Context

## Preferences

- Agents not only gather information and identify competing alternatives, they form **preferences** regarding those alternatives
- Central to **decision theoretic** approaches to rational choice, as applied in philosophy, AI and economics (e.g. Ramsey 1931; Newell 1992)
- Of relevance here: the preferences being mutually taken for granted for the purposes of an interaction
  - Parallel to Stalnaker's common ground

# Preference, Rationality & Context

## Preferences

- A body of preferences can be represented as a binary **preference relation** on the alternatives
- I.e. a set of **pairs of propositions** constructed from  $c$

### Preference State ( $R$ )

- $R$ : binary relation on alternatives (open propositions)
- $R(a, a')$ :  $a$  is preferred to  $a'$
- Each pair in  $R$  is called a *preference*
- Set of (non-empty) alternatives over which  $R$  is defined: issues at stake in  $R$ ,  $C_R$
- Set of worlds among those alternatives: the contextual possibilities written  $c_R$

# Preference, Rationality & Context

## Information in a Preference State

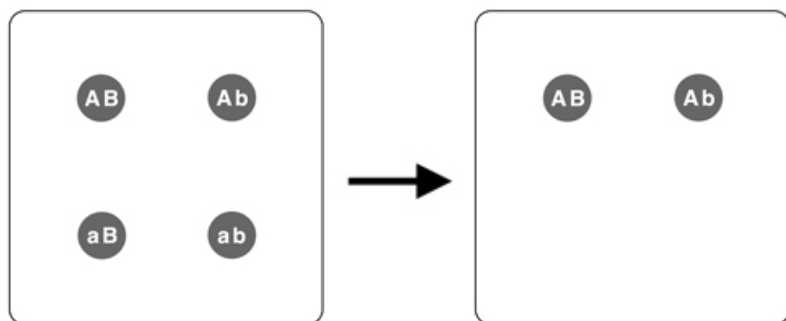


Figure: Accepting the information that A

- $\{w_{AB}, w_{Ab}, w_{aB}, w_{ab}\} \Rightarrow \{w_{AB}, w_{Ab}\}$
- $\{\{\{w_{AB}, w_{Ab}, w_{aB}, w_{ab}\}, \emptyset\}\} \Rightarrow \{\{\{w_{AB}, w_{Ab}\}, \emptyset\}\}$

# Preference, Rationality & Context

## Issues in a Preference State

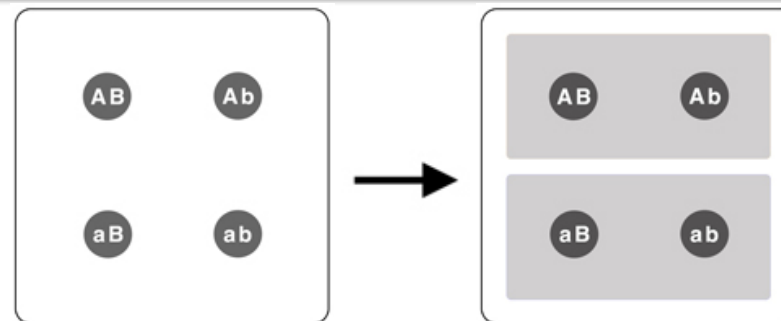
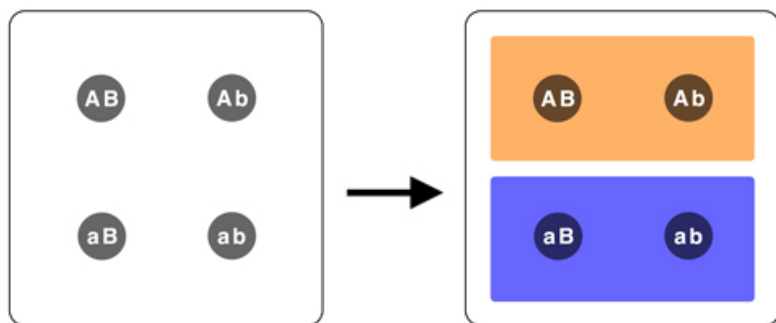


Figure: Recognizing the issue whether A

- $\{\{w_{AB}, w_{Ab}, w_{aB}, w_{ab}\}\} \Rightarrow \{\{w_{AB}, w_{Ab}\}, \{w_{aB}, w_{ab}\}\}$
- $\{\{\{w_{AB}, w_{Ab}, w_{aB}, w_{ab}\}, \emptyset\}\} \Rightarrow \{\{\{w_{AB}, w_{Ab}\}, \emptyset\}, \{\{w_{aB}, w_{ab}\}, \emptyset\}\}$

## Preference, Rationality &amp; Context

Preference and Inquiry

Figure: Coming to prefer A (to  $\neg A$ )

- $\{\{\{w_{AB}, w_{Ab}, w_{aB}, w_{ab}\}, \emptyset\}\}$   
 $\Rightarrow \{\{\{w_{AB}, w_{Ab}\}, \{w_{aB}, w_{ab}\}\}\}$

## Preference, Rationality &amp; Context

Preference and Inquiry

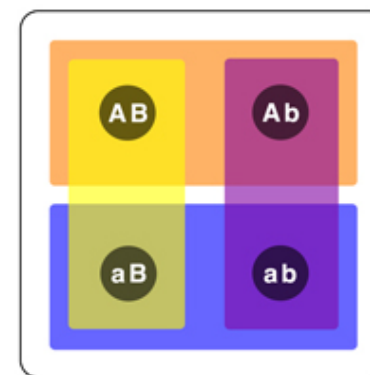


Figure: A more complex preference involving A and B

- $\{\{\{w_{AB}, w_{Ab}\}, \{w_{aB}, w_{ab}\}\}, \{\{w_{AB}, w_{aB}\}, \{w_{Ab}, w_{ab}\}\}\}$

## Preference, Rationality &amp; Context

Using Preference to Make Rational Choices

- Given preference relation, which alternatives are best?
- How do you use preferences to decide what to do?
- In decision theory, this takes the form of defining a **choice function** (Hansson & Grüne-Yanoff 2009)
- A choice function  $Ch$  maps a preference state  $R$  to the set of best alternatives according to  $R$

## Proposal: Choice, Permission, Requirement

- 1  $Ch(R)$  are the alternatives **permissible** according to  $R$
- 2 **Required** by  $R$ : unique alternative permitted by  $R$ 
  - Not always the case!

## Preference, Rationality &amp; Context

The Choice Function: Logical Weak Dominance

## Which Alternatives are Best?

- 1 **Competition between preferred alternatives**  $P(R)$ 
  - Left member in some pair
- 2 If preferred alternative  $a$  is entailed another preferred one, then  $a$  is out
- 3 If  $a$  entails a dispreferred alternative,  $a$  is out

## Choice: Formally

$$Ch(R) = \{a \in P(R) \mid \nexists a' \in P(R) : a' \subset a \\ \& \nexists a' \in D(R) : a \subseteq a'\}$$

$[D(R)$ : dispreferred alternatives]

# Preference, Rationality & Context

How Choice Works: An Example

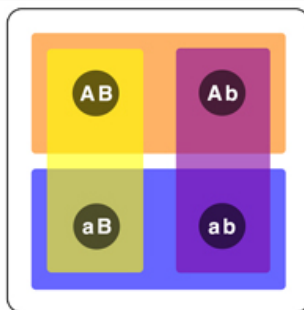


Figure: Preference for A with (separate) preference for B

- $\{\langle \{w_{AB}, w_{Ab}\}, \{w_{aB}, w_{ab}\} \rangle, \langle \{w_{AB}, w_{aB}\}, \{w_{Ab}, w_{ab}\} \rangle\}$
- Two **preferred** (warm) alternatives, orange and yellow
- Neither entails the other nor dispreferred (cold) alt.
- So  $Ch(R) = \{\langle w_{AB}, w_{Ab} \rangle, \langle w_{AB}, w_{aB} \rangle\}$

# Preference, Rationality & Context

How Choice Works: A More Complex Example

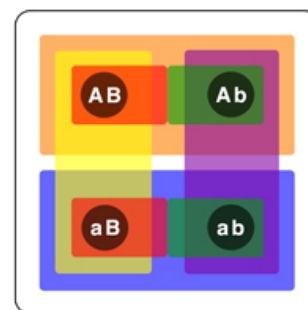


Figure: Pref A and B

$$\{\langle \langle w_{AB}, w_{Ab}, w_{aB}, w_{ab} \rangle, \emptyset \rangle, \langle \langle w_{AB}, w_{Ab} \rangle, \langle w_{aB}, w_{ab} \rangle \rangle, \langle \langle w_{AB} \rangle, \langle w_{Ab} \rangle \rangle, \langle \langle w_{aB} \rangle, \langle w_{ab} \rangle \rangle, \langle \langle w_{AB}, w_{aB} \rangle, \langle w_{Ab}, w_{ab} \rangle \rangle\}$$

- 4 pref. alt's: yellow, orange, reds
- Yellow is out: reds entail it
- Orange is out: top red entails it
- Bottom red is out: it entails blue, which is a dispreferred alt
- Unique best alternative: top red
- $A \wedge B$  is required

# The Semantics: in brief

The New Approach

- **Dynamic Meaning**: function from contents to contents
  - Now contents are preference states
- $R[\phi] = R'$ :  $R'$  is the result of applying  $\phi$  to  $R$

## The Basics

- 1 **May**  $\phi$  tests that  $\phi$  is consistent w/some  $a \in Ch(R)$ 
  - If so, a preference for  $\phi$  is added to  $R$ 
    - This doesn't **guarantee** that  $\phi$  will become  $Ch(R)$
  - Otherwise: fail, i.e. return  $\{\langle \emptyset, \emptyset \rangle\}$
- 2 Disjunction unions separate updates
- 3 Conjunction sequences updates

# The Semantics Applied

How *May* Works: An Example

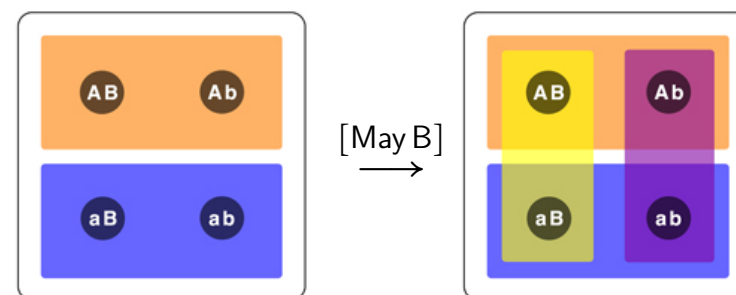


Figure:  $R_0$  to  $R_0[\text{May } B]$

- The best alternative (orange) is consistent w/B
- A preference and permission for B are introduced
- But a requirement for B is **not**



# The Semantics Applied

How *May* Works: With Disjunction

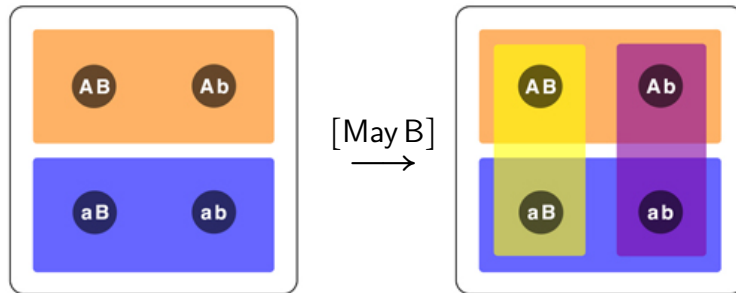


Figure:  $R_0$  to  $R_0[\text{May B}]$

- $R_0$  would be the result of  $R_0[\text{May A}]$
- $R_0[\text{May A} \vee \text{May B}] = R_0[\text{May A}] \cup R_0[\text{May B}] = R_0[\text{May B}]$
- May A  $\not\models$  May A  $\vee$  May B? May A  $\vee$  May B  $\models$  May A?

# Consequence

The Old and the New

## Support, Consequence (Veltman)

- $c \models \phi \iff c[\phi] = c$
- $\phi_1, \dots, \phi_n \models \psi \iff \forall c : c[\phi_1] \dots [\phi_n] \models \psi$

## Preferential Support, Consequence (Starr)

- $R \models \phi \iff Ch(R) = Ch(R[\phi])$
- $\phi_1, \dots, \phi_n \models \psi \iff \forall R : R[\phi_1] \dots [\phi_n] \models \psi$

- Both kinds of consequence and support are useful
- The first when tracking information
- The second when tracking the best alternatives

# The First Problem of Free Choice

Modal Disjunction Intro is Invalid!

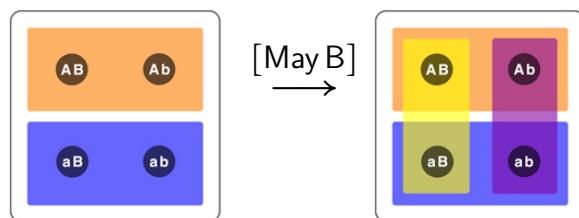


Figure:  $R_0$  to  $R_0[\text{May B}] = R_0[\text{May A} \vee \text{May B}]$

- May A  $\models$  May A  $\vee$  May B means:  
 $Ch(R_0[\text{May A}]) = Ch(R_0[\text{May A}][\text{May A} \vee \text{May B}])$   
 $Ch(R_0) = Ch(R_0[\text{May B}])$
- But yellow is in  $Ch(R_0[\text{May B}])$ , and not  $Ch(R_0)$
- So May A  $\not\models$  May A  $\vee$  May B!

# The Second Problem of Free Choice

From Disjunction to Conjunction

- Disjunctions are felicitous only when each disjunct compatible w/context (Stalnaker 1975: §III)
  - If  $\phi \vee \psi$  is appropriate:  $R[\phi] \neq \{\{\emptyset, \emptyset\}\} \neq R[\psi]$
- What happens when  $\phi$  is May A?
- If May A is compatible w/context then context passes its test and *may* adds preference for A
  - So  $R[\text{May A}]$  will be  $R$  plus a preference for A
- Same holds if  $\psi$  is May B
- Since  $R[\text{May A} \vee \text{May B}] = R[\text{May A}] \cup R[\text{May B}]$ :
  - Resulting state will have pref for A and for B and both will be consistent with best alternatives
- So after updating with May A  $\vee$  May B, May A will not change the best alternatives

# The Second Problem of Free Choice

From Disjunction to Conjunction

## Important Observation

After updating with  $\text{May } A \vee \text{May } B$ ,  $\text{May } A$  will not change the best alternatives

- The same holds for  $\text{May } B$
- But that is just to say:
  - $\text{May } A \vee \text{May } B \models \text{May } A$
  - $\text{May } A \vee \text{May } B \models \text{May } B$
  - $\text{May } A \vee \text{May } B \models \text{May } A \wedge \text{May } B$
- Yet in the modal free fragment,  $\vee$  and  $\wedge$  behave exactly as the Boolean operators of classical semantics!
- **Proviso:** Stalnaker's observation treated as presupposition of  $\vee$  and  $\models$  is Strawsonian

# The Semantics

In Full Detail

## Atomic Semantics

- Where  $R = \{\langle a_0, a_1 \rangle, \dots, \langle a_n, a_{n+1} \rangle\}$ :  

$$R[A] = \{\langle a_0[A], a_1[A] \rangle, \dots, \langle a_n[A], a_{n+1}[A] \rangle\}$$

## Connective Semantics

- $R[\phi \wedge \psi] = (R[\phi])[\psi]$
- $R[\phi \vee \psi] = \begin{cases} R[\phi] \cup R[\psi] & \text{if } R[\phi] \neq \{\langle \emptyset, \emptyset \rangle\} \neq R[\psi] \\ \text{Undefined} & \text{otherwise} \end{cases}$

# The Semantics

And Strawsonian Consequence

## Modal Semantics

$$R[\text{May } \phi] = \begin{cases} R \cup \{\langle c_R[\phi], c_R - c_R[\phi] \rangle\} & \text{if } \exists a \in Ch(R) : a[\phi] \neq \emptyset \\ \{\langle \emptyset, \emptyset \rangle\} & \text{otherwise} \end{cases}$$

$$R[\text{Must } \phi] = \begin{cases} R \cup \{\langle c_R[\phi], c_R - c_R[\phi] \rangle\} & \text{if } \forall a \in Ch(R) : a[\phi] = a \\ \{\langle \emptyset, \emptyset \rangle\} & \text{otherwise} \end{cases}$$

## Strawsonian Preferential Support, Consequence

- $R \models \phi \iff R[\phi]$  is defined &  $Ch(R) = Ch(R[\phi])$
- $\phi_1, \dots, \phi_n \models \psi \iff \forall R : R[\phi_1] \dots [\phi_n] \models \psi$  if  $R[\phi_1] \dots [\phi_n][\psi]$  is defined

# In Related Work

Imperatives

- Free choice behavior of *might*?
  - Veltman's semantics + this implementation of Stalnaker's observation
- Why not just use Veltman's semantics for *may*?
  - *might* has no 'deontic reading'
  - Preference semantics captures deontic reading of *may*
  - Epistemic reading of *may*?
- In other work I use dynamic preference semantics to analyze imperatives
  - Having an analysis of deontic modals allows me to chart their connections
  - E.g.  $!A \models \text{May } A$

# Conclusions

## About Dynamic Semantics

### Summary

- 1 Classical semantics is a fragment of dynamic semantics
- 2 But expanding into the resources afforded only by dynamic semantics is fruitful
  - Here: free choice
- 3 These provide something like an expressivist semantics
  - Truth is not the only useful evaluative concept
  - Consequence is not always about truth
- 4 But avoids typical pitfalls
  - How do you blend the truth-conditional and the non-truth-conditional? (Frege-Geach)
  - DS: truth-conditional behavior emerges for limited fragment from more basic semantic constraints

# Thank you!

(Slides available at <http://williamstarr.net/research>)

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